

# Mixed Sharing Rules

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# The problem

- Each of  $n$  players supplies labour to a joint enterprise under a predetermined rule for sharing the resulting output.
- What are the implications of the chosen sharing rule for (a) efficiency, and (b) distribution, in the resulting noncooperative labour supply game?
- Can a mechanism be devised that leads to an efficient noncooperative equilibrium?

# Assumptions Throughout

- $n$  players
- Technology: neoclassical, diminishing returns to 'labour'  
 $Y = F(L) : F(0) = 0, F'(L) > 0, F''(L) < 0$

- Preferences: well-behaved in every way –  
labour supply irksome

$$u_i = u_i(c_i, l_i) :$$

$$u_{ic}(c_i, l_i) > 0, u_{il}(c_i, l_i) < 0, i = 1, \dots, n$$

Preferences exhibit convexity and **normality**

# Proportional surplus sharing

Assume output shared in proportion to input supplied:  $c_i = \frac{l_i}{L} Y$

$$\Rightarrow u_i = u_i\left(\frac{l_i}{L} F(L), l_i\right) = u_i\left(\frac{l_i}{l_i + L_{-i}} F(l_i + L_{-i}), l_i\right)$$

Player i chooses  $l_i$  to maximise utility:

$$MRS_i = MRT_i$$

or

$$-\frac{\partial u_i(\cdot) / \partial l_i}{\partial u_i(\cdot) / \partial c_i} = \left[\frac{l_i}{L}\right] F'(L) + \left[\frac{L - l_i}{L}\right] \frac{F(L)}{L} = MRT$$

$\Rightarrow$  Overexploitation:  $l_i$  'too high' [Tragedy of the commons]

# Exogenous surplus sharing

$$c_i = \theta_i Y, \sum_{j=1}^n \theta_j = 1, \theta_j \text{'s exogenous}$$

$$\Rightarrow u_i = u_i(\theta_i F(L), l_i) = u_i(\theta_i F(l_i + L_{-i}), l_i)$$

Player  $i$  chooses  $l_i$  to maximise utility:  $MRS_i = MRT_i$

or

$$-\frac{\partial u_i(\cdot) / \partial l_i}{\partial u_i(\cdot) / \partial c_i} = \theta F'(L) < MRT$$

$\Rightarrow$  Under-exploitation:  $l_i$  'too low' [Public good provision]

# A mixed sharing rule

- A sharing rule is initially agreed.
- The rule divides total output into two piles:  $\lambda Y = \lambda F(L)$  is allocated to pile P, and  $(1 - \lambda)Y = [1 - \lambda]F(L)$  to pile E.
- We call the parameter  $\lambda$  the **mixing parameter**.
- The pile P will be divided up proportionally, and the pile E exogenously:

$$c_i = \frac{\ell_i}{L} \lambda F(L) + \theta_i [1 - \lambda] F(L)$$

- Players, knowing the rule, choose their labour inputs noncooperatively

# Questions

1. Does a pure strategy equilibrium exist?
2. If so, is it unique, or are there many?
3. Is there a value of  $\lambda$  which implies a Pareto efficient equilibrium outcome?
4. Can any Pareto efficient outcome be an eq outcome given suitable choices of  $\lambda$  and the  $\theta_i$ 's?
5. Is there a mechanism for selecting a  $\lambda$  that produces a Pareto efficient outcome?

# Method of Analysis

- Define for every player a share function:  $\frac{\hat{\ell}_i}{L} = s_i(L)$
- Observe that, at a Nash equilibrium,

$$\hat{\ell}_1 + \hat{\ell}_2 + \dots + \hat{\ell}_n = L$$

$$\Rightarrow \frac{\hat{\ell}_1}{L} + \frac{\hat{\ell}_2}{L} + \dots + \frac{\hat{\ell}_n}{L} = 1$$

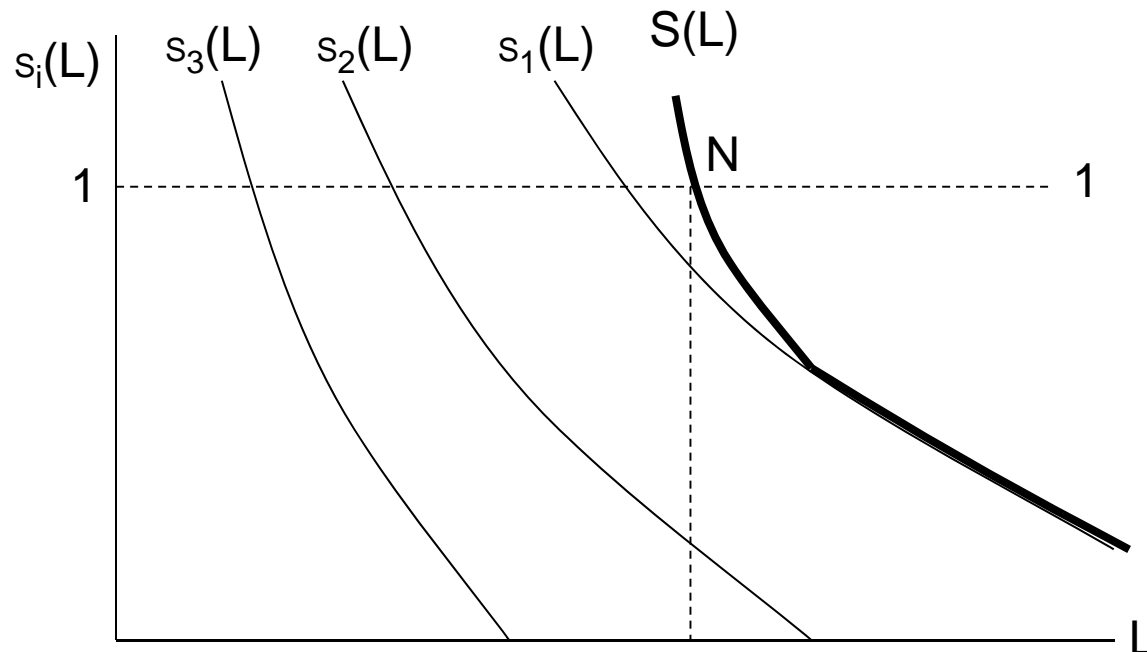
$$\Rightarrow s_1(L) + s_2(L) + \dots + s_n(L) = 1, \text{ or } S(L) = 1$$

- Analysis of Nash equilibrium now involves the solution of one equation in one unknown. [cf the best response function approach, involving n equations.]



# Example

- 3 players,  $\lambda=1/2$ ,  $\theta_1=\theta_2=\theta_3=1/3$
- $U_i() = a_i x_i - \ell_i$ ,  $(a_1, a_2, a_3) = (30, 20, 15)$
- $F(L) = L^{1/2}$



# Some answers

- For a given set of values  $(\lambda, \theta_1, \dots, \theta_n)$ , there exists a unique noncooperative equilibrium.
- Assume identical preferences and equal shares [ $\theta_i=1/n$  for all  $i$ ]. If the mixing parameter  $\lambda$  equals the equilibrium elasticity of production,  $\eta(L)$ , the equilibrium allocation is efficient.

# Another result

- Assume preferences are quasilinear in income. Consider an efficient allocation in which player  $i$  receives output  $x_i^e$  and aggregate input is  $L^e$ . Then the exogenous shares can be chosen so that the equilibrium of the surplus sharing game with  $\lambda = \eta(L^e)$  satisfies  $x_i = x_i^e$  for all  $i$  and  $\sum_i \ell_i^e = L^e$ .

# Freeness from (average) envy

- If the mixing parameter is chosen to equal the elasticity of production, every player prefers her equilibrium bundle to the bundle consisting of  $[F(L^N)/n, L^N/n]$ .

# Unanimity Test

- If  $F(L) = L^\alpha$  and the mixing parameter  $\lambda = \alpha$ , then for every player ,

$$u_i(c_i^N, \ell_i^N) \geq \underset{L \geq 0}{\text{Max}} u_i\left(\frac{F(L)}{n}, \frac{L}{n}\right)$$

Furthermore

$$u_i(c_i^N, \ell_i^N) \geq \underset{L \geq 0}{\text{Max}} u_i(\theta_i F(L), \theta_i L)$$

# The stand-alone test

The stand-alone test formalizes the idea that no player should benefit from the negative externality they impose on others.

Call  $i$  a net contributor if, in equilibrium,

$$l_i/L \geq \theta_i$$

- Then at equilibrium all net contributors pass the stand-alone test.

# Results for large games

Our results may be strengthened for large games  
– that is, games with large  $n$ .

For example:

- The sets of efficient allocations respecting voluntary participation and of equilibrium allocations with optimal mixing are identical.
- By varying the exogenous weights, the whole set of efficient allocations can be mapped out.

# Results for large games (II)

- If exogenous weights are equal, asymptotic equilibria are envy free and pass the unanimity test.
- If exogenous weights are equal, a stronger stand-alone test is satisfied.
- Eq. payoffs for each type are proportional to the same function of the mixing parameter.
- Consequently, all types of player prefer the same value of the mixing parameter.
- **Voting for the optimal value is a dominant strategy for every player in the first stage.**